

## Lectures 16 -17

### QUANTUM SEARCH ALGORITHM (Grover's search)

Suppose that you have  $N$  possible routes to get from one place to another and you would like to find the shortest routes.

Solution: check through all the routes and find the shortest one.

Classical computer requires  $O(N)$  operations to find the shortest way.

Quantum computer requires only  $\sqrt{N}$  operations using Grover's search algorithm.

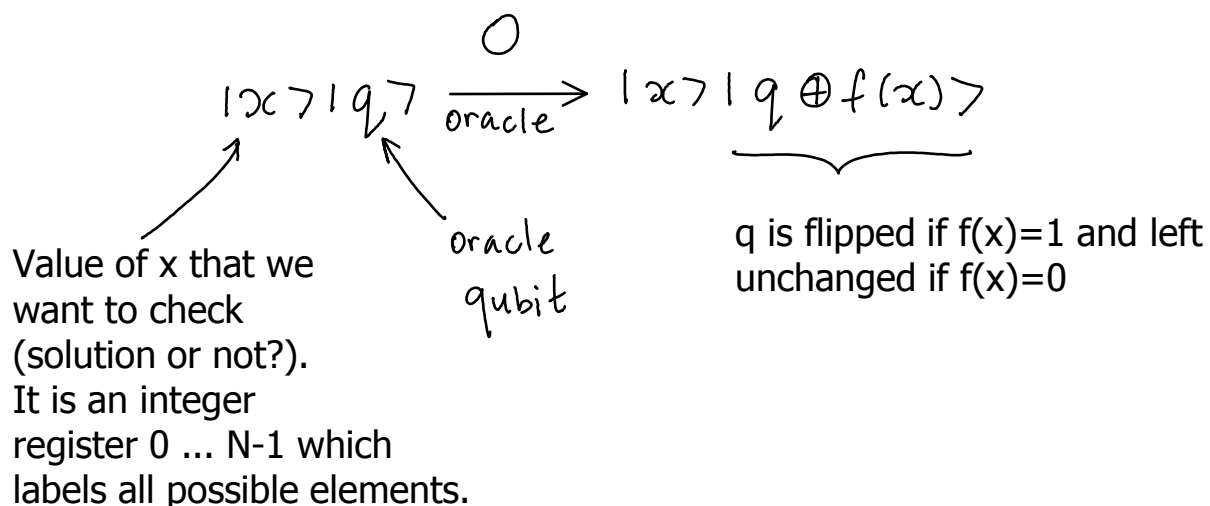
**Problem:** we search through the space of  $N$  elements. Let's deal with the index of the elements:  $0, 1, \dots, N-1$ . We assume for convenience that  $N=2^n$ , i.e. that index can be stored in  $n$  bits. Our search problem has  $M$  solutions:  $1 \leq M \leq N$ .

We define a function  $f(x)$ :

$f(x)=1$  if  $x=0..N-1$  is a solution to our problem

$f(x)=0$  if  $x$  is not a solution.

Now we introduce a **quantum oracle**. It is a black box that can recognize the solutions to the search problem defined above. We will discuss what circuit can be in the black box for a particular example of the search problem later. For now, it is only important what the quantum oracle does.



How to check the solution?

$$|x\rangle |0\rangle \xrightarrow{O} \begin{cases} |x\rangle |0\rangle \\ \text{or} \\ |x\rangle |1\rangle \end{cases} \leftarrow \text{Index } x \text{ corresponds to the element which is a solution to the problem.}$$

Let's change it so the oracle qubit itself does not change.

$$|x\rangle \left[ \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \right] \xrightarrow{O} \begin{cases} \text{not solution} \\ |x\rangle \left[ \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \right] \\ \text{solution} \\ |x\rangle \left[ \frac{1}{\sqrt{2}} (|1\rangle - |0\rangle) \right] \end{cases}$$

$$|x\rangle \left( \frac{|0\rangle - |1\rangle}{\sqrt{2}} \right) \xrightarrow{O} (-1)^{f(x)} |x\rangle \left( \frac{|0\rangle - |1\rangle}{\sqrt{2}} \right)$$

oracle qubit

Remember,  $f(x) = 1$  if  $x$  is a solution  
and  $f(x) = 0$  if  $x$  is not a solution

Oracle qubit is always unchanged  
now so we can omit it from the discussion.

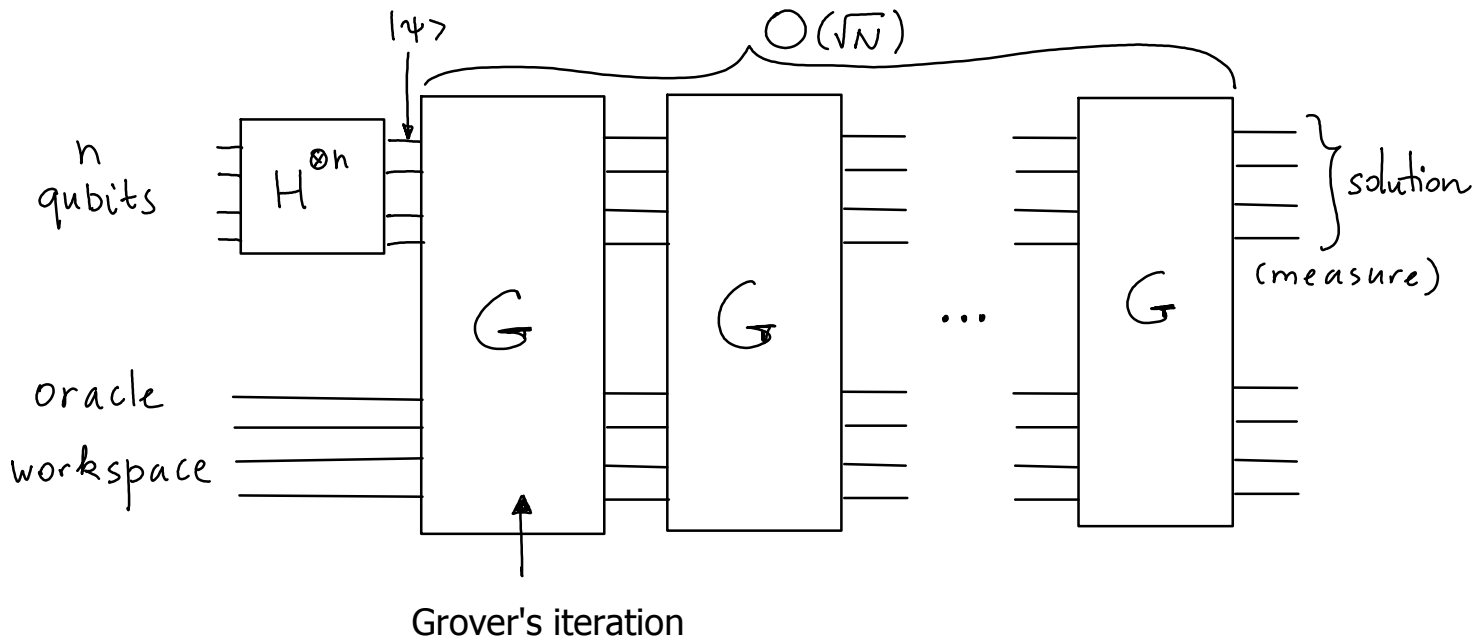
Oracle marks the solution  $|x\rangle \xrightarrow{O} (-1)^{f(x)} |x\rangle$

**Example:** we can factor number  $m$  by checking through all prime numbers from  $x=2$  to  $\sqrt{m}$ . Oracle will calculate  $m/x$  to check if  $x$  is a factor and flip the oracle qubit if it is so. Note: this is not an efficient way to factor.

**Summary: oracle recognizes the solution.**

## Grover iteration & search procedure

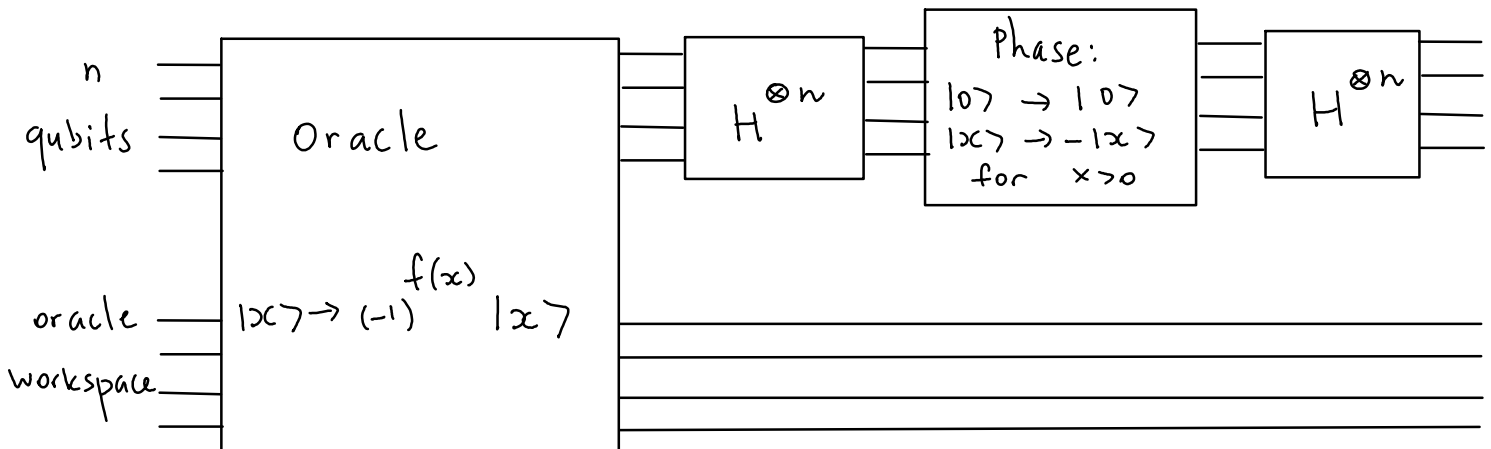
**Goal:** find a solution with least applications of the oracle.



Initial state of the  $N$  qubits :  $|0\rangle^{\otimes n}$

After  $H^{\otimes n}$  :  $|\psi\rangle = \frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} |x\rangle$  (Register is randomized).

### Grover's iteration circuit:



(1) Apply the oracle

(2) Apply the  $H^{\otimes n}$

(3) Conditionally shift phase

(4) Apply the  $H^{\otimes n}$  again

Let's consider step #3 (conditional phase shift) in more detail.  
 State  $|0\rangle$  is the only state which phase is not shifted.

Operator for step 3 is:  $S_3 = 2|0\rangle\langle 0| - I$

Why?

Check its action on  $|x\rangle$ :

If  $|x\rangle \equiv |0\rangle \Rightarrow S_3 |0\rangle = (2|0\rangle\langle 0| - I) |0\rangle = |0\rangle$

If  $|x\rangle \neq |0\rangle \Rightarrow$

$S_3 |x\rangle = (2|0\rangle\langle 0| - I) |x\rangle = -|x\rangle \Rightarrow$

$S_3$  operator shifts phase of  $|x\rangle$  if  $|x\rangle \neq |0\rangle$

$S_2 S_3 S_4$  operator:  $H^{\otimes n} (2|0\rangle\langle 0| - I) H^{\otimes n} =$

$= 2|\psi\rangle\langle\psi| - I$

Remember that  $|\psi\rangle = \frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} |x\rangle$

Therefore, the result of Grover's iteration is:

$G = (2|\psi\rangle\langle\psi| - I) O$  oracle

## What does the Grover iteration do?

We define (normalized) states

$$|\alpha\rangle = \frac{1}{\sqrt{N-M}} \sum_x'' |\alpha\rangle$$

number of elements  $\swarrow$   $\nwarrow$  number of solutions

$\sum''$  indicates sum over  $x$  which are NOT solutions to the problem

and

$$|\beta\rangle = \frac{1}{\sqrt{M}} \sum_x' |\beta\rangle$$

$\nwarrow$  indicates sum over solutions.

Initial state  $|\psi\rangle$ :

$$|\psi\rangle = \sqrt{\frac{N-M}{N}} |\alpha\rangle + \sqrt{\frac{M}{N}} |\beta\rangle = \sum_x'' \sqrt{\frac{N-M}{N}} \frac{1}{\sqrt{N-M}} |\alpha\rangle$$

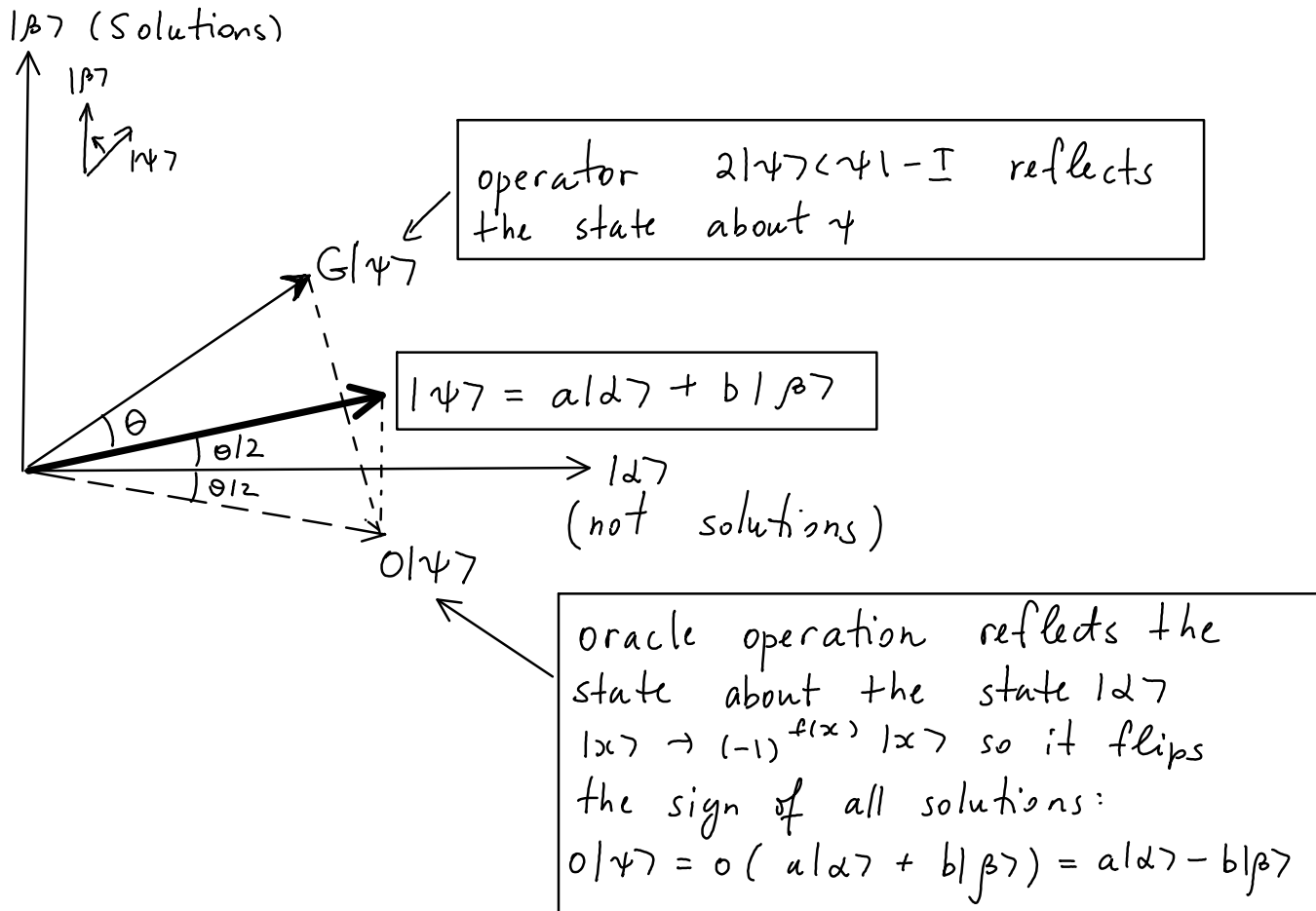
(not solutions)

$$+ \sum_x' \frac{1}{\sqrt{M}} \frac{\sqrt{M}}{\sqrt{N}} |\beta\rangle = \frac{1}{\sqrt{N}} \sum_x |\psi\rangle$$

$\nwarrow$  sum over all states from 0 to  $N-1$ .

Remember: our states  $|\alpha\rangle$  represent indexes of elements 0 ...  $N-1$  to be searched.

## The action of a Grover iteration

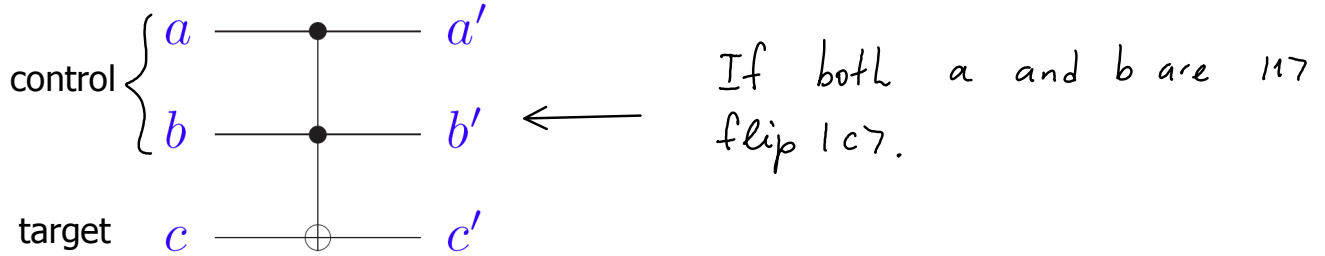


Product of two reflections is a rotation. Therefore, repeated applications of Grover iteration move vector  $|\psi\rangle$  closer to  $|\beta\rangle$ . The measurement will give a solution with high probability since  $|\beta\rangle$  includes all solutions.

# Quantum search: a two-bit example

N = 4

We use a version of Toffoli gate as an oracle.



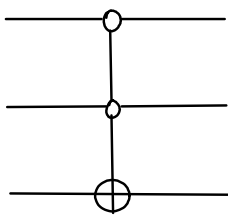
$\frac{107-117}{\sqrt{2}}$ 

 This will work as oracle to mark solution  $x_0 = 3$  since it will change sign for 1117

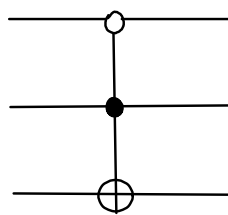
$1117 \frac{107-117}{\sqrt{2}} \rightarrow 1117 \frac{117-107}{\sqrt{2}} = -1117 \frac{107-117}{\sqrt{2}}$

Oracle:  $|x\rangle \xrightarrow{0} (-1)^{f(x)} |x\rangle \rightarrow$  marks the solution  $x_0 = 3$   
 $f(x) = 0$  for  $x_0 = 0, 1, 2$ ;  $f(x) = 1$  for  $x_0 = 3$

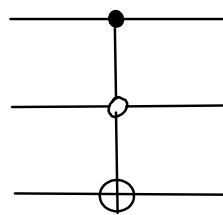
The following versions of Toffoli gate can be used for  $x_0 = 0, 1, 2$ :



$x_0 = 0$   
1007

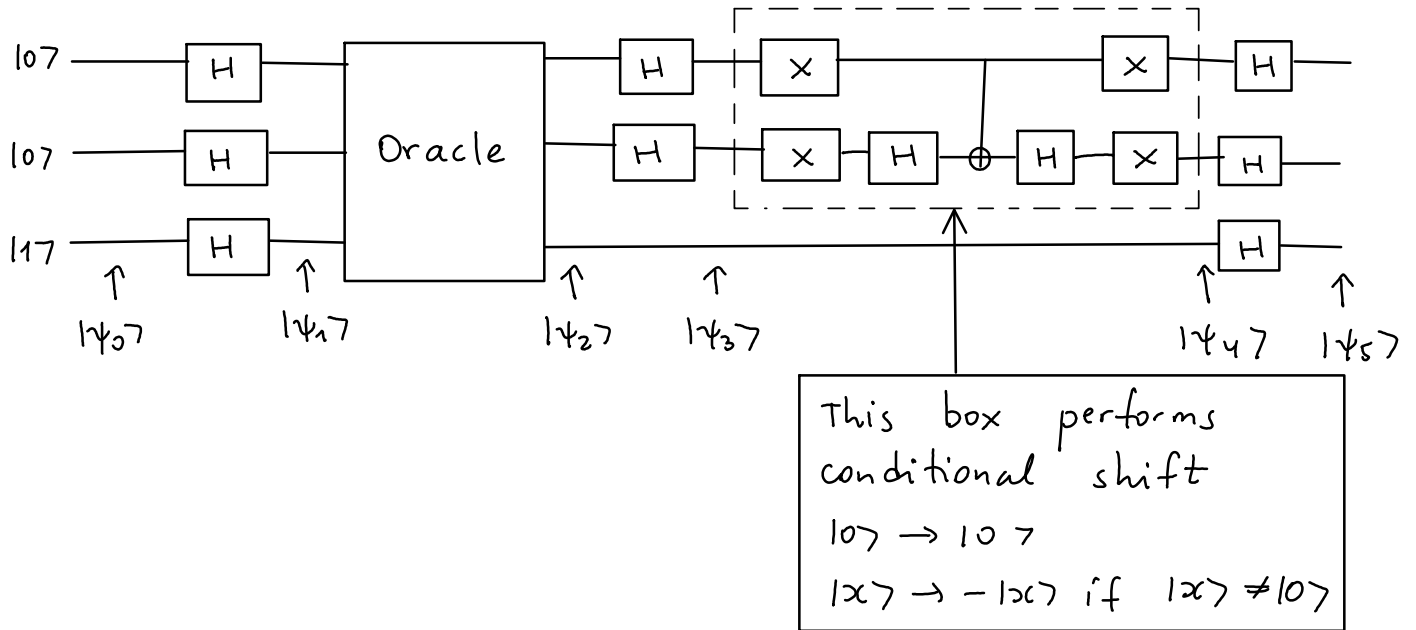


$x_0 = 1$   
1017



$x_0 = 2$   
1107

## Circuit for a two-bit quantum search



**Exercise for the class:** demonstrate that the measurement on first two qubits after this circuit will give  $|01\rangle$  when the corresponding oracle ( $x_0=1$ ) is used.

Our initial state is  $|\psi_0\rangle = |100\rangle$   
↑  
oracle qubit

$$\begin{aligned}
 |\psi_1\rangle &= H|0\rangle H|0\rangle H|1\rangle = \frac{1}{\sqrt{2}} (|0\rangle + |1\rangle) \frac{1}{\sqrt{2}} (|0\rangle + |1\rangle) \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle) \\
 &= \frac{1}{\sqrt{4}} (|00\rangle + |01\rangle + |10\rangle + |11\rangle) \underbrace{\frac{1}{\sqrt{2}} (|0\rangle - |1\rangle)}_{\text{oracle qubit}}
 \end{aligned}$$

$$|\psi_2\rangle = \underbrace{O}_{\text{oracle}} |\psi_1\rangle = \frac{1}{\sqrt{4}} (|00\rangle + |10\rangle + |11\rangle) \frac{|0\rangle - |1\rangle}{\sqrt{2}} + |01\rangle \frac{|1\rangle - |0\rangle}{\sqrt{2}}$$

This version of Toffoli gate

flips  $c$  if  $|ab\rangle = |01\rangle$ .  
 Otherwise, nothing changes.



$$|\psi_2\rangle = \frac{1}{2} [ |00\rangle - |01\rangle + |10\rangle + |11\rangle ] \underbrace{\left[ \frac{|0\rangle - |1\rangle}{\sqrt{2}} \right]}$$

Oracle qubit does not change and is not used in the remaining circuit. Therefore, we can omit it from now on.

Next, we consider how  $H^{\otimes 2}$  gates affect  $|\psi_2\rangle = \frac{1}{2} [ |00\rangle - |01\rangle + |10\rangle + |11\rangle ]$   
(omitted oracle qubit)

$$H^{\otimes 2} |00\rangle = \frac{1}{2} \{ |00\rangle + |01\rangle + |10\rangle + |11\rangle \}$$

$$H^{\otimes 2} |01\rangle = \frac{1}{2} \{ |00\rangle - |01\rangle + |10\rangle - |11\rangle \}$$

$$H^{\otimes 2} |10\rangle = \frac{1}{2} \{ |00\rangle + |01\rangle - |10\rangle - |11\rangle \}$$

$$H^{\otimes 2} |11\rangle = \frac{1}{2} \{ |00\rangle - |01\rangle - |10\rangle + |11\rangle \}$$

$$\begin{aligned} |\psi_3\rangle = H^{\otimes 2} |\psi_2\rangle &= \frac{1}{4} \{ |00\rangle + |01\rangle + |10\rangle + |11\rangle \\ &\quad - |00\rangle + |01\rangle - |10\rangle + |11\rangle \\ &\quad + |00\rangle + |01\rangle - |10\rangle - |11\rangle \\ &\quad + |00\rangle - |01\rangle - |10\rangle + |11\rangle \} \\ &= \frac{1}{2} \{ |00\rangle + |01\rangle - |10\rangle + |11\rangle \} \end{aligned}$$

$$|\psi_4\rangle = S_3 |\psi_3\rangle = \frac{1}{2} \{ |00\rangle - |01\rangle + |10\rangle - |11\rangle \}$$

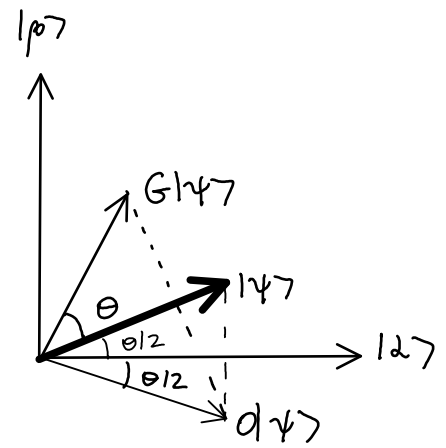
Conditional phase shift, all signs are flipped except for  $|00\rangle$ .

$$\begin{aligned} |\psi_5\rangle = H^{\otimes 2} |\psi_4\rangle &= \frac{1}{4} \{ |00\rangle + |01\rangle + |10\rangle + |11\rangle \\ &\quad - |00\rangle + |01\rangle - |10\rangle + |11\rangle \\ &\quad + |00\rangle + |01\rangle - |10\rangle - |11\rangle \\ &\quad - |00\rangle + |01\rangle + |10\rangle - |11\rangle \} \\ &= \frac{1}{4} \cdot 4 |01\rangle = |01\rangle \equiv |x_0\rangle! \end{aligned}$$



Angle  $\theta$  is determined from:

$$|\psi\rangle = \underbrace{\sqrt{\frac{N-M}{N}}}_{\cos \frac{\theta}{2}} |\alpha\rangle + \underbrace{\sqrt{\frac{M}{N}}}_{\sin \frac{\theta}{2}} |\beta\rangle$$



$$|\psi\rangle = \cos \frac{\theta}{2} |\alpha\rangle + \sin \frac{\theta}{2} |\beta\rangle$$

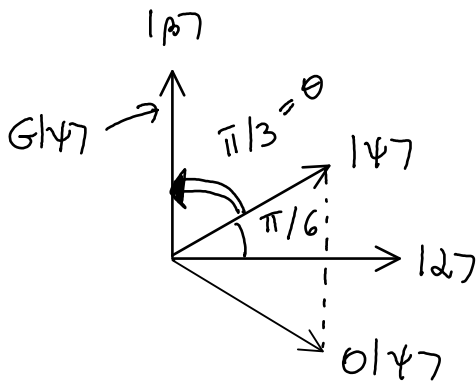
$$0|\psi\rangle = \cos \frac{\theta}{2} |\alpha\rangle - \sin \frac{\theta}{2} |\beta\rangle$$

$$G|\psi\rangle = \cos \frac{3\theta}{2} |\alpha\rangle - \sin \frac{3\theta}{2} |\beta\rangle$$

The  $\theta$  is the rotation angle for Grover iteration.

In our case,

$$|\psi\rangle = \frac{\sqrt{3}}{\sqrt{4}} |\alpha\rangle + \frac{1}{\sqrt{4}} |\beta\rangle \Rightarrow \cos \frac{\theta}{2} = \frac{\sqrt{3}}{2} \quad \boxed{\theta = \frac{\pi}{3}}$$



$$\theta/2 = \pi/6$$

Therefore, one Grover iteration will rotate  $|\psi\rangle$  to  $|\beta\rangle$  exactly.

$$G|\psi\rangle = \underbrace{\cos \frac{3\theta}{2}}_{\cos \frac{\pi}{2} = 0} |\alpha\rangle + \underbrace{\sin \frac{3\theta}{2}}_{\sin \frac{\pi}{2} = 1} |\beta\rangle \equiv |\beta\rangle$$

$$\cos \frac{3\theta}{2} = \cos \frac{\pi}{2} = 0$$

$$\sin \frac{\pi}{2} = 1$$

$$\theta = \pi/3$$