

## Lecture 10

### Quantum error correction

Classical error correction

Modern computers: failure rate is below one error in  $10^{17}$  operations  
Data transmission and storage (file transfers, cell phones, DVD, CD, etc.):  
noise problems

**Solution:** use error-correcting codes to protect against the effect of noise.

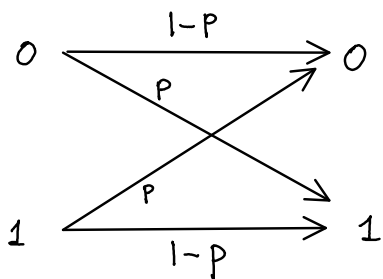
**Key idea:** encode the message with redundant information. Redundancy in the encoded message allows to recover the information in the original message.

### **Classical error correction: repetition code**

**Example:** sending one bit of information across noisy channel.

Effects of the noise: flip the bit with probability  $p$ .

Binary symmetric channel:



Solution: protect the information by making three copy of the bit:

$0 \rightarrow 000$

$1 \rightarrow 111$

Decoding: "majority voting"

011

0 1

vote

1 : 2

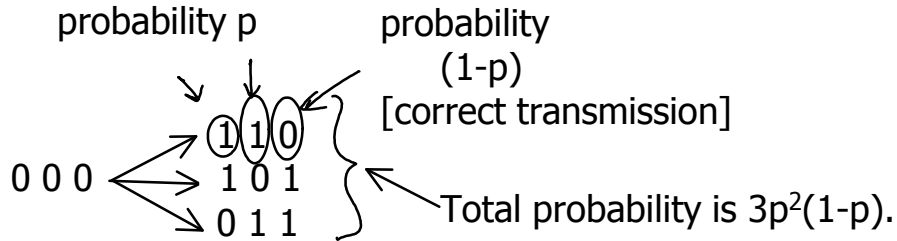
therefore, flip 0 back to 1: 001  $\rightarrow$  111

Note: in this scheme it is impossible to recover the information correctly if more than one bit is flipped.

Probability of error in the non-encoded message:  $p$ .

**Question for the class:** what is the probability of error in the three-bit repetition code if the probability of the bit flip is  $p$ ?

(1) Probability that two bits are flipped:



(2) Probability that three bits are flipped is  $p^3$ :  $000 \longrightarrow 111$  ( $p p p$ ).

Total probability of error  $p_e$  is  $p_e = 3p^2(1-p) + p^3 = 3p^2 - 2p^3$

**Question for the class:** under what circumstances three-bit repetition code is less reliable than the original unencoded transmission?

$$p_e > p \text{ if } 3p^2 - 2p^3 > p \longrightarrow p > 1/2$$

Therefore, three-bit repetition code is less reliable than original one-bit transmission if error rate for a single bit flip exceeds  $1/2$ .

Many clever ideas have been developed for error-correcting codes, always involving redundancy.

## Quantum error correction

### Major difficulties:

(1) No cloning: we can not just duplicate qubit three or more times as in the classical case.

(2) Errors are continuous:

A continuum of different errors may occur on a single qubit. Therefore, determination of which error occurred requires infinite precision and infinite resources.

(3) Measurement collapses the wave function; recovery is impossible if quantum superposition is destroyed.

None of these problems is fatal and all issues may be overcome.

### The three qubit bit flip code

Example: sending the qubit through a channel which flips qubit with probability  $p$ . The state  $|\psi\rangle$  is changed to state  $X|\psi\rangle$ .

Reminder: X gate switches  $|0\rangle$  and  $|1\rangle$

$$X : \alpha|0\rangle + \beta|1\rangle \xrightarrow{X} \alpha|1\rangle + \beta|0\rangle$$
$$X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \quad X \begin{pmatrix} \alpha \\ \beta \end{pmatrix} = \begin{pmatrix} \beta \\ \alpha \end{pmatrix}$$

$|\psi\rangle \xrightarrow{p} X|\psi\rangle \quad (\text{bit flip})$

How does three bit flip code work?

We encode the qubit  $\alpha|0\rangle + \beta|1\rangle$  in three qubits as follows:

$$\alpha|0\rangle + \beta|1\rangle \rightarrow \alpha|000\rangle + \beta|111\rangle$$

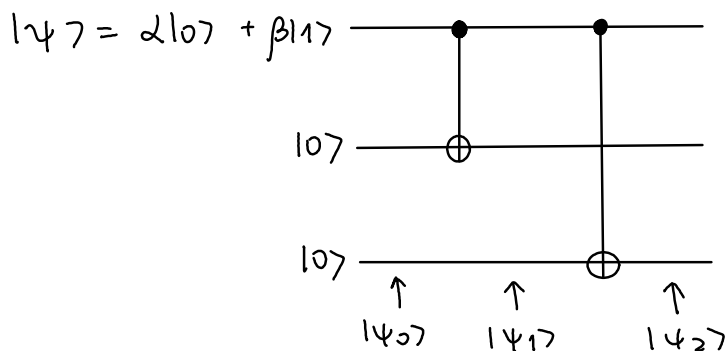
$$|0\rangle \rightarrow |0\rangle_L \equiv |000\rangle$$

$$|1\rangle \rightarrow |1\rangle_L \equiv |111\rangle$$

↑  
"Logical qubits"

Superpositions of basis states are taken to corresponding superpositions of encoded states.

**Exercise to the class:** demonstrate that the logic circuit below implements three bit flip encoding.



Reminder: CNOT gate

$100\rangle \rightarrow 100\rangle$
$101\rangle \rightarrow 101\rangle$
$110\rangle \rightarrow 111\rangle$
$111\rangle \rightarrow 110\rangle$

(flip control qubit if target qubit is  $|1\rangle$ )

$|\psi_0\rangle = \alpha|1000\rangle + \beta|1100\rangle$   
 CNOT: flip 10  $\rightarrow$  11

$|\psi_1\rangle = \alpha|1000\rangle + \beta|1110\rangle$   
 second CNOT: flip 10  $\rightarrow$  11

$|\psi_2\rangle = \alpha|1000\rangle + \beta|1111\rangle$

How does this three-bit flip code works?

Suppose bit flip occurs on one or less qubits. Two-step error correcting procedure is used. Note: each of the qubits is passed through the independent bit-flip channel.

**Error syndrome diagnosis** (or error-detection): we perform the measurement telling us which error, if any, occurred on the quantum state. The measurement result is called the error syndrome.

For the bit-flip channel, there are four error syndromes corresponding to four projection operators:

$$\begin{aligned}
 P_0 &\equiv |000\rangle\langle 000| + |111\rangle\langle 111| && \text{no error} \\
 P_1 &\equiv |100\rangle\langle 100| + |011\rangle\langle 011| && \text{bit flip on qubit \#1} \\
 P_2 &\equiv |010\rangle\langle 010| + |101\rangle\langle 101| && \text{bit flip on qubit \#2} \\
 P_3 &\equiv |001\rangle\langle 001| + |110\rangle\langle 110| && \text{bit flip on qubit \#3}
 \end{aligned}$$

Suppose the bit flip occurred on qubit #1.

Then, the corrupted state is

$$|\psi\rangle = \alpha|100\rangle + \beta|011\rangle.$$

Note that  $\langle\psi|P_1|\psi\rangle = 1$

$$\begin{aligned}
 \text{Proof: } P_1|\psi\rangle &= [ |100\rangle\langle 100| + |011\rangle\langle 011| ] (\alpha|100\rangle + \beta|011\rangle) \\
 &= \alpha|100\rangle + \beta|011\rangle \\
 \langle\psi|P_1|\psi\rangle &= (\langle 100|\alpha^* + \langle 011|\beta^*) (\alpha|100\rangle + \beta|011\rangle) \\
 &= \alpha^*\alpha + \beta^*\beta = 1
 \end{aligned}$$

State of the system after measurement is  $\frac{P_m|\psi\rangle}{\sqrt{p(m)}}$  ;  $p(m) = \langle\psi|P_m|\psi\rangle$

$$m=1 \Rightarrow$$

It is the same as before the measurement  $P_1|\psi\rangle = |\psi\rangle = \alpha|100\rangle + \beta|011\rangle.$

Therefore, the result of the measurement is 1 with 100% probability and the measurement **does not change the state**. We now have the information about which error occurred but we learned nothing about  $\alpha$  or  $\beta$ .

**Recovery:** the value of the error syndrome is used to recover the initial state. The four possible error syndromes and the recovery procedures in each case are:

- 0 (no error) - do nothing
- 1 (bit flip on the first qubit) - flip the first qubit again (use X gate)
- 2 (bit flip on the second qubit) - flip the second qubit again
- 3 (bit flip on the third qubit) - flip the third qubit again

### Exercise for the class:

What is the corrupted state  $|\psi\rangle$  if the bit flip occurred on qubit #3?

Show that  $p(3)=1$  and  $P_3|\psi\rangle=|\psi\rangle$ , so that the detection of this error by measurement will not change the state of the system.

$$|\psi\rangle = \alpha|001\rangle + \beta|110\rangle$$

$$P_3|\psi\rangle = |001\rangle\langle 001|(\alpha|001\rangle + \beta|110\rangle) + |110\rangle\langle 110|(\alpha|001\rangle + \beta|110\rangle) = \alpha|001\rangle + \beta|110\rangle$$

$$p(3) = \langle\psi|P_3|\psi\rangle = (\langle 001|\alpha^* + \langle 110|\beta^*)$$

$$(|001\rangle\alpha + |110\rangle\beta) = \alpha^*\alpha + \beta^*\beta = 1$$